

CMPT 409/981: Optimization for Machine Learning

Lecture 15

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Online Optimization

- 1: Online Optimization (w_0 , Algorithm \mathcal{A} , Convex set \mathcal{C})
 - 2: **for** $k = 1, \dots, T$ **do**
 - 3: Algorithm \mathcal{A} chooses point (decision) $w_k \in \mathcal{C}$
 - 4: Environment chooses and reveals the (potentially adversarial) loss function $f_k : \mathcal{C} \rightarrow \mathbb{R}$
 - 5: Algorithm suffers a cost $f_k(w_k)$
 - 6: **end for**
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Regret: For any fixed decision $u \in \mathcal{C}$, $R_T(u) := \sum_{k=1}^T [f_k(w_k) - f_k(u)]$.

Online Gradient Descent (OGD): At iteration k , OGD chooses w_k . After the loss function f_k is revealed, OGD uses the function to compute

$$w_{k+1} = \Pi_C[w_k - \eta_k \nabla f_k(w_k)] \text{ where } \Pi_C[x] = \arg \min_{y \in \mathcal{C}} \frac{1}{2} \|y - x\|^2 .$$

If the convex set \mathcal{C} has a diameter D i.e. for all $x, y \in \mathcal{C}$, $\|x - y\| \leq D$, for an arbitrary sequence losses such that each f_k is convex, differentiable and G -Lipschitz, OGD with $\eta_k = \frac{D}{\sqrt{2} G \sqrt{k}}$ and $w_1 \in \mathcal{C}$, has regret $R_T(u) \leq \sqrt{2} DG \sqrt{T}$.

Additionally, if each f_k is μ_k strongly-convex, OGD with $\eta_k = \frac{1}{\sum_{i=1}^k \mu_i}$ has regret $R_T(u) \leq \frac{G^2}{2\mu} (1 + \log(T))$.

Follow the Leader (FTL): At iteration k , FTL chooses the point w_k . After the loss function f_k is revealed, FTL uses it to compute

$$w_{k+1} = \arg \min_{w \in \mathcal{C}} \sum_{i=1}^k f_i(w).$$

Running FTL on a quadratic lower-bound for the loss recovers OGD in the strongly-convex case.

For strongly-convex, G -Lipschitz losses, FTL has regret $R_T(u) \leq \frac{G^2}{2\mu} (1 + \log(T))$ that matches OGD, but does not require knowledge of μ (Proof today).

If the losses are not necessarily strongly-convex, then FTL can result in $O(T)$ regret.

Recap

Idea: Add an explicit regularization to fix FTL for a convex sequence of losses.

Follow the Regularized Leader (FTRL): At iteration $k \geq 0$, FTRL chooses the point w_k . After the loss function f_k is revealed, FTRL uses it to compute

$$w_{k+1} = \arg \min_{w \in \mathcal{C}} \sum_{i=1}^k \left[f_i(w) + \frac{\sigma_i}{2} \|w - w_i\|^2 \right] + \frac{\sigma_0}{2} \|w\|^2,$$

where $\sigma_i \geq 0$ is the regularization strength. If we set $\sigma_i = 0$ for all i , FTRL reduces to FTL.

Running FTRL on a linear lower-bound for the loss recovers OGD in the convex case.

FTRL has the following regret for a general sequence of convex losses,

$$R_T(u) \leq \sum_{k=1}^T \left[\frac{1}{2\lambda_{k+1}} \|\nabla f_k(w_k)\|^2 \right] + \sum_{k=1}^T \frac{\sigma_k}{2} \|u - w_k\|^2 + \frac{\sigma_0}{2} \|u\|^2 \quad \text{where } \lambda_k = \sum_{i=1}^{k-1} [\mu_i] + \sum_{i=0}^k [\sigma_i].$$

For convex, G -Lipschitz losses, FTRL has regret $R_T(u) \leq \sqrt{2} \sqrt{D^2 + \|u\|^2} G \sqrt{T}$.

Follow the Leader - Strongly-Convex, Lipschitz functions

Claim: If the convex set \mathcal{C} has diameter D , for an arbitrary sequence losses such that each f_k is μ_k strongly-convex (s.t. $\mu := \min_{k=1}^T \mu_k > 0$), G -Lipschitz and differentiable, then FTL with $w_1 \in \mathcal{C}$ satisfies the following regret bound for all $u \in \mathcal{C}$,

$$R_T(u) \leq \frac{G^2}{2\mu} (1 + \log(T))$$

Proof: Using the general result for FTRL, for $\lambda_{k+1} = \sum_{i=1}^k \mu_i + \sum_{i=0}^k \sigma_i$. Since f_k is μ_k strongly-convex, we will set $\sigma_i = 0$ for all i . Hence, $\lambda_{k+1} = \sum_{i=1}^k \mu_i \geq \mu k$.

$$R_T(u) \leq \sum_{k=1}^T \left[\frac{1}{2\lambda_{k+1}} \|\nabla f_k(w_k)\|^2 \right] + \sum_{i=1}^T \frac{\sigma_i}{2} \|u - w_i\|^2 + \frac{\sigma_0}{2} \|u\|^2 \leq \frac{G^2}{2\mu} \sum_{k=1}^T \left[\frac{1}{k} \right]$$

(Since f_k is G -Lipschitz)

$$\implies R_T(u) \leq \frac{G^2 (1 + \log(T))}{2\mu}$$

Hence, FTL matches the regret for OGD for strongly-convex, Lipschitz functions, but does not require knowledge of μ .

Questions?

Adaptive step-sizes

Recall the claim we proved in Lecture 14 (Slide 6): If the convex set \mathcal{C} has diameter D , for an arbitrary sequence of losses such that each f_k is convex and differentiable, OGD with the update $w_{k+1} = \Pi_{\mathcal{C}}[w_k - \eta_k \nabla f_k(w_k)]$ such that $\eta_k \leq \eta_{k-1}$ and $w_1 \in \mathcal{C}$ has the following regret for $u \in \mathcal{C}$,

$$R_T(u) \leq \frac{D^2}{2\eta_T} + \sum_{k=1}^T \frac{\eta_k}{2} \|\nabla f_k(w_k)\|^2 = \frac{D^2}{2\eta} + \frac{\eta}{2} \sum_{k=1}^T \|\nabla f_k(w_k)\|^2 \quad (\text{If } \eta_k = \eta \text{ for all } k)$$

In order to find the optimal η , differentiating the RHS w.r.t η and setting it to zero,

$$-\frac{D^2}{2\eta^2} + \frac{1}{2} \sum_{k=1}^T \|\nabla f_k(w_k)\|^2 = 0 \implies \eta^* = \frac{D}{\sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2}}$$

Since the second derivative equal to $\frac{2D^2}{\eta^3} > 0$, η^* minimizes the RHS. Setting $\eta = \eta^*$,

$$R_T(u) \leq D \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2}$$

Adaptive step-sizes

Choosing $\eta = \eta^* = \frac{D}{\sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2}}$ minimizes the upper-bound on the regret. However, this is not practical since setting η requires knowing $\nabla f_k(w_k)$ for all $k \in [T]$.

To approximate η^* to have a practical algorithm, we can set η_k as follows:

$$\eta_k = \frac{D}{\sqrt{\sum_{s=1}^k \|\nabla f_s(w_s)\|^2}}$$

Hence, at iteration k , we only use the gradients upto that iteration.

Algorithmically, we only need to maintain the running sum of the squared gradient norms.

Moreover, this choice of step-size ensures that $\eta_k \leq \eta_{k-1}$ (since we are accumulating gradient norms in the denominator so the step-size cannot increase) and hence we can use our general result for bounding the regret.

Scalar AdaGrad

Hence, we have the following update for any $\eta > 0$,

$$w_{k+1} = \Pi_C[w_k - \eta_k \nabla f_k(w_k)] \quad ; \quad \eta_k = \frac{\eta}{\sqrt{\sum_{s=1}^k \|\nabla f_s(w_s)\|^2}}$$

This is exactly the AdaGrad update without a per-coordinate scaling and is referred to as scalar AdaGrad or AdaGrad Norm [WWB20].

For a sequence of convex, differentiable losses, using the general result,

$$R_T(u) \leq \frac{D^2}{2\eta_T} + \sum_{k=1}^T \frac{\eta_k}{2} \|\nabla f_k(w_k)\|^2 = \frac{D^2}{2\eta} \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2} + \frac{\eta}{2} \sum_{k=1}^T \frac{\|\nabla f_k(w_k)\|^2}{\sqrt{\sum_{s=1}^k \|\nabla f_s(w_s)\|^2}}$$

In order to bound the regret for AdaGrad, we need to bound the last term.

Scalar AdaGrad

We prove the following general claim and will use it for $a_s = \|\nabla f_s(w_s)\|^2$.

Claim: For all T and $a_s \geq 0$, $\sum_{k=1}^T \frac{a_k}{\sqrt{\sum_{s=1}^k a_s}} \leq 2\sqrt{\sum_{k=1}^T a_k}$.

Proof: Let us prove by induction. **Base case:** For $T = 1$, LHS = $\sqrt{a_1} < 2\sqrt{a_1} =$ RHS.

Inductive Hypothesis: If the statement is true for $T - 1$, we need to prove it for T .

$$\sum_{k=1}^T \frac{a_k}{\sqrt{\sum_{s=1}^k a_s}} = \sum_{k=1}^{T-1} \frac{a_k}{\sqrt{\sum_{s=1}^k a_s}} + \frac{a_T}{\sqrt{\sum_{s=1}^T a_s}} \leq 2\sqrt{\sum_{s=1}^{T-1} a_s} + \frac{a_T}{\sqrt{\sum_{s=1}^T a_s}} = 2\sqrt{Z-x} + \frac{x}{\sqrt{Z}}$$

$(x := a_T, Z := \sum_{s=1}^T a_s)$

The derivative of the RHS w.r.t to x is $-\frac{1}{\sqrt{Z-x}} + \frac{1}{\sqrt{Z}} < 0$ for all $x \geq 0$ and hence the RHS is maximized at $x = 0$. Setting $x = 0$ completes the induction proof.

$$\Rightarrow \sum_{k=1}^T \frac{a_k}{\sqrt{\sum_{s=1}^k a_s}} \leq 2\sqrt{Z} = 2\sqrt{\sum_{s=1}^T a_s}$$

Scalar AdaGrad

Recall that $R_T(u) \leq \frac{D^2}{2\eta} \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2} + \frac{\eta}{2} \sum_{k=1}^T \frac{\|\nabla f_k(w_k)\|^2}{\sqrt{\sum_{s=1}^k \|\nabla f_s(w_s)\|^2}}$. Using the claim in the previous slide with $a_s := \|\nabla f_s(w_s)\|^2 \geq 0$,

$$R_T(u) \leq \frac{D^2}{2\eta} \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2} + \eta \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2} = \left(\frac{D^2}{2\eta} + \eta \right) \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2}.$$

The step-size that minimizes the above bound is equal to $\eta^* = \frac{D}{\sqrt{2}}$. With this choice,

$$R_T(u) \leq \sqrt{2}D \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2}$$

Comparing to the regret for the optimal (impractical) constant step-size on Slide 3,

$$R_T(u) \leq \sqrt{2} \min_{\eta} \left[\frac{D^2}{2\eta} + \frac{\eta}{2} \sum_{k=1}^T \|\nabla f_k(w_k)\|^2 \right]$$

Hence, AdaGrad is only sub-optimal by $\sqrt{2}$ when compared to the best constant step-size!

Scalar AdaGrad - Convex, Lipschitz functions

Claim: If the convex set \mathcal{C} has diameter D i.e. for all $x, y \in \mathcal{C}$, $\|x - y\| \leq D$, for an arbitrary sequence losses such that each f_k is convex, differentiable and G -Lipschitz, scalar AdaGrad with $\eta_k = \frac{\eta}{\sqrt{\sum_{s=1}^k \|\nabla f_s(w_s)\|^2}}$ and $w_1 \in \mathcal{C}$ has the following regret for all $u \in \mathcal{C}$,

$$R_T(u) \leq \left(\frac{D^2}{2\eta} + \eta \right) G \sqrt{T}$$

Proof: Using the general result from the previous slide,

$$R_T(u) \leq \left(\frac{D^2}{2\eta} + \eta \right) \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2} \leq \left(\frac{D^2}{2\eta} + \eta \right) \sqrt{G^2 T} = \left(\frac{D^2}{2\eta} + \eta \right) G \sqrt{T}$$

(Since each f_k is G -Lipschitz)

With $\eta = \frac{D}{\sqrt{2}}$, $R_T(u) \leq \sqrt{2} D G \sqrt{T}$. Hence, for convex, Lipschitz functions, AdaGrad achieves the same regret as OGD but is adaptive to G .

Scalar AdaGrad - Strongly-Convex, Lipschitz functions

Claim: If the convex set \mathcal{C} has diameter D i.e. for all $x, y \in \mathcal{C}$, $\|x - y\| \leq D$, for an arbitrary sequence losses such that each f_k is μ strongly-convex, differentiable and G -Lipschitz, scalar AdaGrad with $\eta_k = \frac{G^2/\mu}{1 + \sum_{s=1}^k \|\nabla f_s(w_s)\|^2}$ and $w_1 \in \mathcal{C}$ has the following regret for all $u \in \mathcal{C}$,

$$R_T(u) \leq \frac{G^2}{2\mu} [1 + \log(1 + G^2 T)]$$

Though AdaGrad can achieve logarithmic regret for strongly-convex, Lipschitz functions similar to OGD and FTL, it requires knowledge of G and μ and is not adaptive to these quantities.

Proof: Need to prove this in Assignment 4!

Questions?

Let us consider a more general and practical variant of AdaGrad that uses a per-coordinate step-size. The corresponding update is:

$$v_{k+1} = w_k - \eta A_k^{-1} \nabla f_k(w_k) \quad ; \quad w_{k+1} = \Pi_C^k[v_{k+1}] := \arg \min_{w \in C} \frac{1}{2} \|w - v_{k+1}\|_{A_k}^2 .$$

$$A_k = \begin{cases} \sqrt{\sum_{s=1}^k \|\nabla f_s(w_s)\|^2} I_d & \text{(Scalar AdaGrad)} \\ \text{diag}(G_k^{\frac{1}{2}}) & \text{(Diagonal AdaGrad)} \\ G_k^{\frac{1}{2}} & \text{(Full-Matrix AdaGrad)} \end{cases}$$

where $G_k \in \mathbb{R}^{d \times d} := \sum_{s=1}^k [\nabla f_s(w_s) \nabla f_s(w_s)^\top]$. For the subsequent analysis, we will assume that A_k is invertible (a small ϵI_d can be added to ensure invertibility)

Claim: If the convex set \mathcal{C} has diameter D , for an arbitrary sequence of losses such that each f_k is convex and differentiable, AdaGrad with the general update $w_{k+1} = \Pi_{\mathcal{C}}^k[w_k - \eta A_k^{-1} \nabla f_k(w_k)]$ and $w_1 \in \mathcal{C}$ has the following regret for $u \in \mathcal{C}$,

$$R_T(u) \leq \left(\frac{D^2}{2\eta} + \eta \right) \sqrt{d} \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2}$$

Proof: Starting from the update, $v_{k+1} = w_k - \eta A_k^{-1} \nabla f_k(w_k)$,

$$v_{k+1} - u = w_k - \eta A_k^{-1} \nabla f_k(w_k) - u \implies A_k[v_{k+1} - u] = A_k[w_k - u] - \eta \nabla f_k(w_k)$$

Multiplying the above equations,

$$\begin{aligned} [v_{k+1} - u]^T A_k [v_{k+1} - u] &= [w_k - u - \eta A_k^{-1} \nabla f_k(w_k)]^T [A_k[w_k - u] - \eta \nabla f_k(w_k)] \\ \|v_{k+1} - u\|_{A_k}^2 &= \|w_k - u\|_{A_k}^2 - 2\eta \langle \nabla f_k(w_k), w_k - u \rangle + \eta^2 [A_k^{-1} \nabla f_k(w_k)]^T [\nabla f_k(w_k)] \\ \implies \|v_{k+1} - u\|_{A_k}^2 &= \|w_k - u\|_{A_k}^2 - 2\eta \langle \nabla f_k(w_k), w_k - u \rangle + \eta^2 \|\nabla f_k(w_k)\|_{A_k^{-1}}^2 \end{aligned}$$

Recall that $\|v_{k+1} - u\|_{A_k}^2 = \|w_k - u\|_{A_k}^2 - 2\eta \langle \nabla f_k(w_k), w_k - u \rangle + \eta^2 \|\nabla f_k(w_k)\|_{A_k^{-1}}^2$. Using the update $w_{k+1} = \Pi_{\mathcal{C}}^k[v_{k+1}]$, $u \in \mathcal{C}$ with the non-expansiveness of projections,

$$\begin{aligned} \|w_{k+1} - u\|_{A_k}^2 &= \|\Pi_{\mathcal{C}}[v_{k+1}] - \Pi_{\mathcal{C}}[u]\|_{A_k}^2 \leq \|v_{k+1} - u\|_{A_k}^2 \\ \implies \|w_{k+1} - u\|_{A_k}^2 &\leq \|w_k - u\|_{A_k}^2 - 2\eta \langle \nabla f_k(w_k), w_k - u \rangle + \eta^2 \|\nabla f_k(w_k)\|_{A_k^{-1}}^2 \\ &\leq \|w_k - u\|_{A_k}^2 - 2\eta [f_k(w_k) - f_k(u)] + \eta^2 \|\nabla f_k(w_k)\|_{A_k^{-1}}^2 \quad (\text{Convexity}) \\ \implies f_k(w_k) - f_k(u) &\leq \frac{\|w_k - u\|_{A_k}^2 - \|w_{k+1} - u\|_{A_k}^2}{2\eta} + \frac{\eta}{2} \|\nabla f_k(w_k)\|_{A_k^{-1}}^2 \end{aligned}$$

Summing from $k = 1$ to T ,

$$\implies R_T(u) \leq \frac{1}{2\eta} \sum_{k=1}^T \left[\|w_k - u\|_{A_k}^2 - \|w_{k+1} - u\|_{A_k}^2 \right] + \frac{\eta}{2} \sum_{k=1}^T \|\nabla f_k(w_k)\|_{A_k^{-1}}^2$$

Let us now bound the first term in the above expression.

$$\begin{aligned}
& \sum_{k=1}^T \left[\|w_k - u\|_{A_k}^2 - \|w_{k+1} - u\|_{A_k}^2 \right] \\
&= \sum_{k=2}^T [(w_k - u)^\top [A_k - A_{k-1}] (w_k - u)] + \|w_1 - u\|_{A_1}^2 - \|w_{T+1} - u\|_{A_T}^2 \\
&\leq \sum_{k=2}^T \|w_k - u\|^2 \lambda_{\max}[A_k - A_{k-1}] + \|w_1 - u\|_{A_1}^2 \leq \sum_{k=2}^T D^2 \lambda_{\max}[A_k - A_{k-1}] + \|w_1 - u\|_{A_1}^2 \\
&\quad \text{(Since } A_{k-1} \preceq A_k, \lambda_{\max}[A_k - A_{k-1}] \geq 0 \text{ and } \|w_k - u\|^2 \leq D) \\
&\implies \sum_{k=1}^T \left[\|w_k - u\|_{A_k}^2 - \|w_{k+1} - u\|_{A_k}^2 \right] \leq D^2 \sum_{k=2}^T \text{Tr}[A_k - A_{k-1}] + \|w_1 - u\|_{A_1}^2 \\
&\quad \text{(For any PSD matrix } B, \lambda_{\max}[B] \leq \text{Tr}[B])
\end{aligned}$$

Continuing the proof from the previous slide,

$$\begin{aligned}
 \sum_{k=1}^T \left[\|w_k - u\|_{A_k}^2 - \|w_{k+1} - u\|_{A_k}^2 \right] &\leq D^2 \sum_{k=2}^T \text{Tr}[A_k - A_{k-1}] + \|w_1 - u\|_{A_1}^2 \\
 &= D^2 \text{Tr} \left[\sum_{k=2}^T [A_k - A_{k-1}] \right] + \|w_1 - u\|_{A_1}^2 && \text{(Linearity of Trace)} \\
 &= D^2 \text{Tr}[A_T - A_1] + \|w_1 - u\|_{A_1}^2 \leq D^2 \text{Tr}[A_T - A_1] + \lambda_{\max}[A_1] \|w_1 - u\|^2 \\
 \sum_{k=1}^T \left[\|w_k - u\|_{A_k}^2 - \|w_{k+1} - u\|_{A_k}^2 \right] &\leq D^2 \text{Tr}[A_T] - D^2 \text{Tr}[A_1] + D^2 \text{Tr}[A_1] = D^2 \text{Tr}[A_T]
 \end{aligned}$$

Putting everything together,

$$R_T(u) \leq \frac{D^2 \text{Tr}[A_T]}{2\eta} + \frac{\eta}{2} \sum_{k=1}^T \|\nabla f_k(w_k)\|_{A_k^{-1}}^2$$

Let us now bound the second term in the above expression.

Claim: $\sum_{k=1}^T \|\nabla f_k(w_k)\|_{A_k^{-1}}^2 \leq 2 \operatorname{Tr}[A_T]$

Proof: Let us prove by induction. For convenience, define $\nabla_k := \nabla f_k(w_k)$.

Base case: For $k = 1$, $\text{LHS} = \operatorname{Tr}[\nabla_1^\top A_1^{-1} \nabla_1] = \operatorname{Tr}[A_1^{-1} \nabla_1 \nabla_1^\top] = \operatorname{Tr}[A_1^{-1} A_1 A_1] \leq 2 \operatorname{Tr}[A_1] = \text{RHS}$. Here, we used the cyclic property of trace i.e. $\operatorname{Tr}[ABC] = \operatorname{Tr}[BCA]$.

Inductive Hypothesis: If the statement is true for $T - 1$, we need to prove it for T .

$$\sum_{k=1}^{T-1} \|\nabla_k\|_{A_k^{-1}}^2 + \|\nabla_T\|_{A_T^{-1}}^2 \leq 2 \operatorname{Tr}[A_{T-1}] + \|\nabla_T\|_{A_T^{-1}}^2 = 2 \operatorname{Tr}[(A_T^2 - \nabla_T \nabla_T^\top)^{1/2}] + \operatorname{Tr}[A_T^{-1} \nabla_T \nabla_T^\top]$$

For any $X \succeq Y \succeq 0$, we have [DHS11, Lemma 8], $2 \operatorname{Tr}[(X - Y)^{1/2}] + \operatorname{Tr}[X^{-1/2} Y] \leq 2 \operatorname{Tr}[X^{1/2}]$.

Using this for $X = A_T^2$, $Y = \nabla_T \nabla_T^\top$, $\sum_{k=1}^T \|\nabla_k\|_{A_k^{-1}}^2 \leq 2 \operatorname{Tr}[A_T]$, which completes the proof.

Putting everything together,



$$R_T(u) \leq \left(\frac{D^2}{2\eta} + \eta \right) \operatorname{Tr}[A_T].$$

Recall that $R_T(u) \leq \left(\frac{D^2}{2\eta} + \eta\right) \text{Tr}[A_T]$. Bounding $\text{Tr}[A_T]$

$$\begin{aligned} \text{Tr}[A_T] &= \text{Tr}[G_T^{\frac{1}{2}}] = \sum_{j=1}^d \sqrt{\lambda_j[G_T]} = d \frac{\sum_{j=1}^d \sqrt{\lambda_j[G_T]}}{d} \leq d \sqrt{\frac{\sum_{j=1}^d \lambda_j[G_T]}{d}} \\ & \hspace{15em} \text{(Jensen's inequality for } \sqrt{x} \text{)} \\ &= \sqrt{d} \sqrt{\sum_{j=1}^d \lambda_j[G_T]} = \sqrt{d} \sqrt{\text{Tr}[G_T]} = \sqrt{d} \sqrt{\text{Tr} \left[\sum_{k=1}^T \nabla f_k(w_k) \nabla f_k(w_k)^\top \right]} \\ \text{Tr}[A_T] &\leq \sqrt{d} \sqrt{\left[\sum_{k=1}^T \text{Tr} \nabla f_k(w_k) \nabla f_k(w_k)^\top \right]} = \sqrt{d} \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2} \quad \text{(Linearity of Trace)} \end{aligned}$$

Putting everything together,

$$R_T(u) \leq \left(\frac{D^2}{2\eta} + \eta\right) \sqrt{d} \sqrt{\sum_{k=1}^T \|\nabla f_k(w_k)\|^2}$$

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-  Rachel Ward, Xiaoxia Wu, and Leon Bottou, *Adagrad stepsizes: Sharp convergence over nonconvex landscapes*, The Journal of Machine Learning Research **21** (2020), no. 1, 9047–9076.