# CMPT 210: Probability and Computing 

Lecture 14

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March 5, 2024

## Recap

Probability density function (PDF): Let $R$ be a r.v. with codomain $V$. The probability density function of $R$ is the function $\mathrm{PDF}_{R}: V \rightarrow[0,1]$, such that $\operatorname{PDF}_{R}[x]=\operatorname{Pr}[R=x]$ if $x \in \operatorname{Range}(\mathrm{R})$ and equal to zero if $x \notin \operatorname{Range}(\mathrm{R})$.
Cumulative distribution function (CDF): The cumulative distribution function of $R$ is the function $\mathrm{CDF}_{R}: \mathbb{R} \rightarrow[0,1]$, such that $\mathrm{CDF}_{R}[x]=\operatorname{Pr}[R \leq x]$.

Importantly, neither $\mathrm{PDF}_{R}$ nor $\mathrm{CDF}_{R}$ involves the sample space of an experiment.
Example: If we flip three coins, and $C$ counts the number of heads, then
$\operatorname{PDF}_{C}[0]=\operatorname{Pr}[C=0]=\frac{1}{8}$, and
$\mathrm{CDF}_{C}[2.3]=\operatorname{Pr}[C \leq 2.3]=\operatorname{Pr}[C=0]+\operatorname{Pr}[C=1]+\operatorname{Pr}[C=2]=\frac{7}{8}$.

## Recap

A distribution can be specified by its probability density function (PDF) (denoted by $f$ ).
Bernoulli Distribution: $f_{p}(0)=1-p, f_{p}(1)=p$. Example: When tossing a coin such that $\operatorname{Pr}[$ heads $]=p$, random variable $R$ is equal to 1 if we get a heads (and equal to 0 otherwise). In this case, $R$ follows the Bernoulli distribution i.e. $R \sim \operatorname{Ber}(p)$.
Uniform Distribution: If $R: \mathcal{S} \rightarrow V$, then for all $v \in V, f(v)=1 /|V|$. Example: When throwing an $n$-sided die, random variable $R$ is the number that comes up on the die. $V=\{1,2, \ldots, n\}$. In this case, $R$ follows the Uniform distribution i.e. $R \sim \operatorname{Uniform}(1, n)$.
Binomial Distribution: $f_{n, p}(k)=\binom{n}{k} p^{k}(1-p)^{n-k}$. Example: When tossing $n$ independent coins such that $\operatorname{Pr}[$ heads $]=p$, random variable $R$ is the number of heads in $n$ coin tosses. In this case, $R$ follows the Binomial distribution i.e. $R \sim \operatorname{Bin}(n, p)$.
Geometric Distribution: $f_{p}(k)=(1-p)^{k-1} p$. Example: When repeatedly tossing a coin such that $\operatorname{Pr}[$ heads $]=p$, random variable $R$ is the number of tosses needed to get the first heads. In this case, $R$ follows the Geometric distribution i.e. $R \sim \operatorname{Geo}(p)$.

## Distributions - Examples

Q: It is known that disks produced by a certain company will be defective with probability 0.01 independently of each other. The company sells the disks in packages of 10 and offers a money-back guarantee that at most 1 of the 10 disks is defective (the package can be returned if there is more than 1 defective disk). What proportion of packages is returned? If someone buys three packages, what is the probability that exactly one of them will be returned?

Let $X$ be the random variable corresponding to the number of defective disks in a package. Let $E$ be the event that the package is returned. We wish to compute $\operatorname{Pr}[E]=\operatorname{Pr}[X>1] . X$ follows the Binomial distribution $\operatorname{Bin}(10,0.01)$. Hence,

$$
\begin{aligned}
\operatorname{Pr}[E]=\operatorname{Pr}[X>1] & =1-\operatorname{Pr}[X \leq 1]=1-\operatorname{Pr}[X=0]-\operatorname{Pr}[X=1] \\
& =1-\binom{10}{0}(0.99)^{10}-\binom{10}{1}(0.99)^{9}(0.01)^{1} \approx 0.05
\end{aligned}
$$

## Distributions - Examples

Q: It is known that disks produced by a certain company will be defective with probability 0.01 independently of each other. The company sells the disks in packages of 10 and offers a money-back guarantee that at most 1 of the 10 disks is defective (the package can be returned if there is more than 1 defective disk). If someone buys three packages, what is the probability that exactly one of them will be returned?

Let $F$ be the event that someone bought 3 packages and exactly one of them is returned.
Answer 1: Let $E_{i}$ be the event that package $i$ is returned. From the previous question, we know that $\operatorname{Pr}\left[E_{i}\right]=\operatorname{Pr}[$ Package $i$ has more than 1 defective disk $] \approx 0.05$.

$$
\begin{aligned}
& F=\left(E_{1} \cap E_{2}^{c} \cap E_{3}^{c}\right) \cup\left(E_{1}^{c} \cap E_{2}^{c} \cap E_{3}\right) \cup\left(E_{1}^{c} \cap E_{2} \cap E_{3}^{c}\right) \\
& \operatorname{Pr}[F]=\operatorname{Pr}\left[E_{1}\right]\left(1-\operatorname{Pr}\left[E_{2}\right]\right)\left(1-\operatorname{Pr}\left[E_{3}\right]\right)+\left(1-\operatorname{Pr}\left[E_{1}\right]\right)\left(1-\operatorname{Pr}\left[E_{2}\right]\right) \operatorname{Pr}\left[E_{3}\right]+\ldots \\
& \operatorname{Pr}[F] \approx 3 \times(0.05)(0.95)(0.95) \approx 0.15
\end{aligned}
$$

Answer 2: Let $Y$ be the random variable corresponding to the number of packages returned. $Y$ follows the Binomial distribution $\operatorname{Bin}(3,0.05)$ and we wish to compute $\operatorname{Pr}[F]=\operatorname{Pr}[Y=1] \approx\binom{3}{1}(0.05)^{1}(0.95)^{2} \approx 0.15$.

## Distributions - Examples

Q: You are randomly and independently throwing darts. The probability that you hit the bullseye in throw $i$ is $p$. Once you hit the bullseye you win and can go collect your reward. (a) What is the probability that you win after exactly $k$ throws? (b) What is the probability you win in less than $k$ throws?
(a) The number of throws $(T)$ to hit the bullseye and win follows a geometric distribution $\operatorname{Geo}(p)$ and we wish to compute $\operatorname{Pr}[T=k]$. Using the PDF for the Geometric distribution, this is equal to $(1-p)^{k-1} p$.
(b) Answer 1: If $E$ is the event that we win in less than $k$ throws, $\operatorname{Pr}[E]=\operatorname{Pr}[T<k]=\sum_{i=1}^{k-1} \operatorname{Pr}[T=i]=p \sum_{i=1}^{k-1}(1-p)^{i-1}=1-(1-p)^{k-1}$.

## Answer 2:

$\operatorname{Pr}[E]=1-\operatorname{Pr}\left[E^{c}\right]=1-\operatorname{Pr}[$ do not hit the bullseye in $k-1$ throws $]=1-(1-p)^{k-1}$.

## Number Guessing Game

Q: We have two envelopes. Each contains a distinct number in $\{0,1,2, \ldots, 100\}$. To win the game, we must determine which envelope contains the larger number. We are allowed to peek at the number in one envelope selected at random. Can we devise a winning strategy?

Strategy 1: We pick an envelope at random and guess that it contains the larger number (without even peeking at the number).

Q: What is the probability that we win with this strategy? Ans: 0.5
Strategy 2: We peek at the number and if its below 50 , we choose the other envelope.
But the numbers in the envelopes need not be random! The numbers are chosen "adversarially" in a way that will defeat our guessing strategy. For example, to "beat" Strategy 2, the two numbers can always be chosen to be below 50 .

Q: Can we do better than $50 \%$ chance of winning?

## Number Guessing Game

Suppose that we somehow knew a number $x$ that was in between the numbers in the envelopes. If we peek in one envelope and see a number. If it is bigger than $x$, we know its the higher number and choose that envelope. If it is smaller than $x$, we know that is the smaller number and choose the other envelope.
Of course, we do not know such a number $x$. But we can guess it!
Strategy 3: Choose a random number $x$ from $\{0.5,1.5,2.5, \ldots n-1 / 2\}$ according to the uniform distribution i.e. $\operatorname{Pr}[x=0.5]=\operatorname{Pr}[1.5]=\ldots=1 / n$. Then we peek at the number (denoted by $T$ ) in one envelope, and if $T>x$, we choose that envelope, else we choose the other envelope.

The advantage of such a randomized strategy is that the adversary cannot easily "adapt" to it. Q: But does it have better than $50 \%$ chance of winning?

## Number Guessing Game

Let the numbers in the two envelopes be $L$ (lower number) and $H$ (the higher number).


$$
\begin{aligned}
\operatorname{Pr}[\text { win }] & =\frac{L}{2 n}+\frac{H-L}{2 n}+\frac{H-L}{2 n}+\frac{n-H}{2 n} \\
& =\frac{1}{2}+\frac{H-L}{2 n} \geq \frac{1}{2}+\frac{1}{2 n}>\frac{1}{2}
\end{aligned}
$$

Hence our strategy has a greater than $50 \%$ chance of winning! If $n=10, \operatorname{Pr}[\mathrm{win}] \geq 0.55$, for $n=100, \operatorname{Pr}[\mathrm{win}] \geq 0.505$.
Q: For $n=100$, if $L=23$ and $H=54$, compute $\operatorname{Pr}[$ guessing too low $\mid$ we win ]
Ans: $\operatorname{Pr}[$ guessing too low $\mid$ we win $]=$ $\frac{\operatorname{Pr}[\text { we win } \cap \text { guessing too low }]}{\operatorname{Pr}[\text { we win }]}=\frac{L / 2 n}{1 / 2+(H-L) / 2 n}=$ $\frac{L}{n+H-L}=\frac{23}{131}$.

## Questions?

