# CMPT 210: Probability and Computing 

Lecture 12

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## Recap - (Basic) Frievald's Algorithm

- Q: For $n \times n$ matrices $A, B$ and $D$, is $D=A B$ ?
- Last class, we proved that:

$$
\left\lvert\, \begin{array}{c|c|c|} 
& \text { Yes } & \text { No } \\
D=A B & 1 & 0 \\
D \neq A B & <\frac{1}{2} & \geq \frac{1}{2}
\end{array}\right.
$$

Table 1: Probabilities for Basic Frievalds Algorithm

## Frievald's Algorithm

By repeating the Basic Frievald's Algorithm $m$ times, we will amplify the probability of success. The resulting complete Frievald's Algorithm is given by:

1 Run the Basic Frievald's Algorithm for $m$ independent runs.
2 If any run of the Basic Frievald's Algorithm outputs "no", output "no".
3 If all runs of the Basic Frievald's Algorithm output "yes", output "yes".

$$
\left\lvert\, \begin{array}{c|c|c|} 
& \text { Yes } & \text { No } \\
D=A B & 1 & 0 \\
D \neq A B & <\frac{1}{2^{m}} & \geq 1-\frac{1}{2^{m}}
\end{array}\right.
$$

Table 2: Probabilities for Frievald's Algorithm

If $m=20$, then Frievald's algorithm will make mistake with probability $1 / 2^{20} \approx 10^{-6}$.
Computational Complexity: $O\left(m n^{2}\right)$

## Probability Amplification

Consider a randomized algorithm $\mathcal{A}$ that is supposed to solve a binary decision problem i.e. it is supposed to answer either Yes or No. It has a one-sided error - (i) if the true answer is Yes, then the algorithm $\mathcal{A}$ correctly outputs Yes with probability 1, but (ii) if the true answer is No, the algorithm $\mathcal{A}$ incorrectly outputs Yes with probability $\leq \frac{1}{2}$.

Let us define a new algorithm $\mathcal{B}$ that runs algorithm $\mathcal{A} m$ times, and if any run of $\mathcal{A}$ outputs No, algorithm $\mathcal{B}$ outputs No. If all runs of $\mathcal{A}$ output Yes, algorithm $\mathcal{B}$ outputs Yes.
Q: What is the probability that algorithm $\mathcal{B}$ correctly outputs Yes if the true answer is Yes, and correctly outputs No if the true answer is No?

## Probability Amplification - Analysis

$$
\begin{aligned}
& \text { If } A_{i} \text { denotes run } i \text { of Algorithm } \mathcal{A} \text {, then } \\
& \\
& \operatorname{Pr}[\mathcal{B} \text { outputs Yes } \mid \text { true answer is Yes }] \\
& =\operatorname{Pr}\left[\mathcal{A}_{1} \text { outputs Yes } \cap \mathcal{A}_{2} \text { outputs Yes } \cap \ldots \cap \mathcal{A}_{m} \text { outputs Yes } \mid \text { true answer is Yes }\right] \\
& =\prod_{i=1}^{m} \operatorname{Pr}\left[\mathcal{A}_{i} \text { outputs Yes } \mid \text { true answer is Yes }\right]=1 \quad \text { (Independence of runs) } \\
& \operatorname{Pr}[\mathcal{B} \text { outputs } \mathrm{No} \mid \text { true answer is No }] \\
& =1-\operatorname{Pr}[\mathcal{B} \text { outputs Yes } \mid \text { true answer is No }] \\
& =1-\operatorname{Pr}\left[\mathcal{A}_{1} \text { outputs Yes } \cap \mathcal{A}_{2} \text { outputs Yes } \cap \ldots \cap \mathcal{A}_{m} \text { outputs Yes } \mid \text { true answer is No }\right] \\
& =1-\prod_{i=1}^{m} \operatorname{Pr}\left[\mathcal{A}_{i} \text { outputs Yes } \mid \text { true answer is } \mathrm{No}\right] \geq 1-\frac{1}{2^{m}} .
\end{aligned}
$$

When the true answer is Yes, both $\mathcal{B}$ and $\mathcal{A}$ correctly output Yes. When the true answer is No, $\mathcal{A}$ incorrectly outputs Yes with probability $<\frac{1}{2}$, but $\mathcal{B}$ incorrectly outputs Yes with probability $<\frac{1}{2^{m}} \ll \frac{1}{2}$. By repeating the experiment, we have "amplified" the probability of success.

## Questions?

## Random Variables

Definition: A random "variable" $R$ on a probability space is a total function whose domain is the sample space $\mathcal{S}$. The codomain is usually a subset of the real numbers.

Example: Suppose we toss three independent, unbiased coins. Let $C$ be the number of heads that appear.
$\mathcal{S}=\{H H H, H H T, H T H$, HTT, THH, THT, TTH, TTT $\}$
$C$ is a total function that maps each outcome in $\mathcal{S}$ to a number as follows: $C(H H H)=3$, $C(H H T)=C(H T H)=C(T H H)=2, C(H T T)=C(T H T)=C(T T H)=1, C(T T T)=0$.
$C$ is a random variable that counts the number of heads in 3 tosses of the coin.
Example: I toss a coin, and define the random variable $R$ which is equal to 1 when I get a heads, and equal to 0 when I get a tails.

Bernoulli random variables: Random variables with the codomain $\{0,1\}$ are called Bernoulli random variables. E.g. $R$ is a Bernoulli r.v.

## Back to throwing dice

Q: Suppose we throw two standard dice one after the other. Let us define $R$ to be the random variable equal to the sum of the dice. What is the domain, range of $R$ ?

Ans: $R:\{1,2,3,4,5,6\} \times\{1,2,3,4,5,6\} \rightarrow \mathbb{N} \cap[2,12]$.
$R((4,7))=11, R((4,1))=5, R((1,1))=2, R((6,6))=12$.
Q: Three balls are randomly selected from an urn containing 20 balls numbered 1 through 20. The random variable $M$ is the maximal value on the selected balls. What is the domain, range of $M$ ? Ans: $M:\{1,2, \ldots, 20\} \times\{1,2, \ldots, 20\} \times\{1,2, \ldots, 20\} \rightarrow\{1,2, \ldots, 20\}$

Q: In the above example, what is $2 \times M((1,4,6))$ ? Is $M$ an invertible function? Ans: 12 , No since $M$ maps both $\{1,2,5)$ and $(3,4,5)$ to 5 .

## Random Variables and Events

Indicator Random Variable: An indicator random variable maps every outcome to either 0 or 1. Example: Suppose we throw two standard dice, and define $M$ to be the random variable that is 1 iff both throws of the dice produce a prime number, else it is 0 .
$M:\{1,2,3,4,5,6\} \times\{1,2,3,4,5,6\} \rightarrow\{0,1\} . M((2,3))=1, M((3,6))=0$.
An indicator random variable partitions the sample space into those outcomes mapped to 1 and those outcomes mapped to 0 .

Example: When throwing two dice, if $E$ is the event that both throws of the dice result in a prime number, then random variable $M=1$ iff event $E$ happens, else $M=0$.

The indicator random variable corresponding to an event $E$ is denoted as $\mathcal{I}_{E}$, meaning that for $\omega \in E, \mathcal{I}_{E}[\omega]=1$ and for $\omega \notin E, \mathcal{I}_{E}[\omega]=0$. In the above example, $M=\mathcal{I}_{E}$ and since $(2,4) \notin E, M((2,4))=0$ and since $(3,5) \in E, M((3,5))=1$.

## Random Variables and Events

In general, a random variable that takes on several values partitions $\mathcal{S}$ into several blocks.
Example: When we toss a coin three times, and define $C$ to be the r.v. that counts the number of heads, $C$ partitions $\mathcal{S}$ as follows: $\mathcal{S}=\{\underbrace{H H H}_{C=3}, \underbrace{H H T, H T H, T H H}_{C=2}, \underbrace{H T T, T H T, T T H}_{C=1}, \underbrace{T T T}_{C=0}\}$.
Each block is a subset of the sample space and is therefore an event. For example, [ $C=2$ ] is the event that the number of heads is two and consists of the outcomes $\{H H T, H T H, T H H\}$.

Since it is an event, we can compute its probability i.e.
$\operatorname{Pr}[C=2]=\operatorname{Pr}[\{H H T, H T H, T H H\}]=\operatorname{Pr}[\{H H T\}]+\operatorname{Pr}[\{H T H\}]+\operatorname{Pr}[\{T H H\}]$. Since this is a uniform probability space, $\operatorname{Pr}[\omega]=\frac{1}{8}$ for $\omega \in \mathcal{S}$ and hence $\operatorname{Pr}[C=2]=\frac{3}{8}$.
Q: What is $\operatorname{Pr}[C=0], \operatorname{Pr}[C=1]$ and $\operatorname{Pr}[C=3]$ ? Ans: $\frac{1}{8}, \frac{3}{8}, \frac{1}{8}$
Q: What is $\sum_{i=0}^{3} \operatorname{Pr}[C=i]$ ? Ans: 1
Since a random variable $R$ is a total function that maps every outcome in $\mathcal{S}$ to some value in the codomain, $\sum_{i \in \text { Range of } R} \operatorname{Pr}[R=i]=\sum_{i \in \text { Range of } R} \sum_{\omega \text { s.t. }} R(\omega)=i \operatorname{Pr}[\omega]=\sum_{\omega \in \mathcal{S}} \operatorname{Pr}[\omega]=1$.

## Back to throwing dice

Q: Suppose we throw two standard dice one after the other. Let us define $R$ to be the random variable equal to the sum of the dice. What are the outcomes in the event $[R=2]$ ? Ans: $\{(1,1)\}$

Q: What is $\operatorname{Pr}[R=4], \operatorname{Pr}[R=9]$ ? Ans: $\frac{3}{36}, \frac{4}{36}$
Q: If $M$ is the indicator random variable equal to 1 iff both throws of the dice produces a prime number, what is $\operatorname{Pr}[M=1]$ ? Ans: $\frac{9}{36}$

## Distribution Functions

Probability density function (PDF): Let $R$ be a random variable with codomain $V$. The probability density function of $R$ is the function $\mathrm{PDF}_{R}: V \rightarrow[0,1]$, such that
$\operatorname{PDF}_{R}[x]=\operatorname{Pr}[R=x]$ if $x \in \operatorname{Range}(\mathrm{R})$ and equal to zero if $x \notin \operatorname{Range}(\mathrm{R})$.
$\sum_{x \in V} \operatorname{PDF}_{R}[x]=\sum_{x \in \operatorname{Range}(R)} \operatorname{Pr}[R=x]=1$.
Cumulative distribution function (CDF): If the codomain is a subset of the real numbers, then the cumulative distribution function is the function $\mathrm{CDF}_{R}: \mathbb{R} \rightarrow[0,1]$, such that $\mathrm{CDF}_{R}[x]=\operatorname{Pr}[R \leq x]$.
Importantly, neither $\mathrm{PDF}_{R}$ nor $\mathrm{CDF}_{R}$ involves the sample space of an experiment.
Example: If we flip three coins, and $C$ counts the number of heads, then
$\operatorname{PDF}_{c}[0]=\operatorname{Pr}[C=0]=\frac{1}{8}$, and
$\mathrm{CDF}_{C}[2.3]=\operatorname{Pr}[C \leq 2.3]=\operatorname{Pr}[C=0]+\operatorname{Pr}[C=1]+\operatorname{Pr}[C=2]=\frac{7}{8}$.
Q: What is CDF $_{C}[5.8]$ ? Ans: 1 .
For a general random variable $R$, as $x \rightarrow \infty, \operatorname{CDF}_{R}[x] \rightarrow 1$ and $x \rightarrow-\infty, \mathrm{CDF}_{R}[x] \rightarrow 0$.

## Back to throwing dice

Q: Suppose we throw two standard dice one after the other. Let us define $T$ to be the random variable equal to the sum of the dice. Plot $\mathrm{PDF}_{T}$ and $\mathrm{CDF}_{T}$

Recall that $T:\{1,2,3,4,5,6\} \times\{1,2,3,4,5,6\} \rightarrow V$ where $V=\{2,3,4, \ldots 12\}$.
$\mathrm{PDF}_{T}: V \rightarrow[0,1]$ and $\mathrm{CDF}_{T}: \mathbb{R} \rightarrow[0,1]$.
For example, $\operatorname{PDF}_{T}[4]=\operatorname{Pr}[T=4]=\frac{3}{36}$ and $\operatorname{PDF}_{T}[12]=\operatorname{Pr}[T=12]=\frac{1}{36}$.



## Questions?

## Distributions

Many random variables turn out to have the same PDF and CDF. In other words, even though $R$ and $T$ might be different random variables on different probability spaces, it is often the case that $\mathrm{PDF}_{R}=\mathrm{PDF}_{T}$. Hence, by studying the properties of such PDFs, we can study different random variables and experiments.

Distribution over a random variable can be fully specified using the cumulative distribution function (CDF) (usually denoted by $F$ ). The corresponding probability density function (PDF) is denoted by $f$.

Common Discrete Distributions in Computer Science:

- Bernoulli Distribution
- Uniform Distribution
- Binomial Distribution
- Geometric Distribution

