## **CMPT 210:** Probability and Computation

Lecture 1

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- Course Webpage: https://vaswanis.github.io/210-S22
- Piazza: https://piazza.com/sfu.ca/summer2022/cmpt210/home
- Prerequisites: MACM 101, MATH 152 and MATH 232/MATH 240

## **Course Information**

**Objective**: Introduce the foundational concepts in probability required by computing.

#### Syllabus:

- Combinatorics, Set Theory, Inclusion-Exclusion.
- Probability theory, Random variables, Joint distributions.
- Expectation, Variance, Standard Deviation, Discrete distributions: Binomial and Geometric.
- Conditional probability, Bayes' Theorem, Tail inequalities (Markov, Chebyshev, Chernoff).
- Applications: Freivalds' algorithm, Quicksort, Max-cut, Load Balancing, A/B testing
- Normal Distribution, Central Limit Theorem (introduction)

#### Resources:

- Introduction to Probability and Statistics for Engineers and Scientists (Ross).
- Mathematics for Computer Science (Meyer, Lehman, Leighton): https://people.csail.mit.edu/meyer/mcs.pdf
- CMU Lecture Notes for Probability and Computing (O'Donnell): http://www.cs.cmu. edu/~odonnell/papers/probability-and-computing-lecture-notes.pdf

### • Grading:

- 4 Assignments (4  $\times$  12.5% = 50%)
- 1 Mid-Terms (1  $\times$  15% = 15%) (24 June)
- 1 Final Exam (1  $\times$  35% = 35%) (TBD)
- Each assignment is due in 1 week (typically Friday).
- For some flexibility, each student is allowed 1 late-submission and can submit in the next class (typically the Tuesday after).
- If you miss the mid-term (needs to be a well-justified reason), will reassign weight to the final.
- If you miss the final, there will be a make-up exam.

# Questions?

**Informal def**: Unordered collection of objects (referred to as *elements*) **Examples**:  $\{a, b, c\}$ ,  $\{\{a, b\}, \{c, a\}\}$ ,  $\{1.2, 2.5\}$ ,  $\{$ yellow, red, green $\}$ ,  $\{x|x \text{ is capital of a North American country}\}$ ,  $\{x|x \text{ is an integer in } [5, 10]\}$ .

There is no notion of an element appearing twice. E.g.  $\{a, a, b\} = \{a, b\}$ .

The order of the elements does not matter. E.g.  $A = \{a, b\} = \{b, a\}$ .

 $C = \{x | x \text{ is a color of the rainbow } \}$ 

**Elements** of C: red, orange, yellow, green, blue, indigo, violet.

**Membership**: red  $\in$  *C*, brown  $\notin$  *C*.

**Cardinality**: Number of elements in the set. |C| = 7

Q: A = {x|5 < x < 17 and x is a power of 2 }. Enumerate A. What is |A|? Ans: A = {8,16}, |A| = 2

### **Common Sets**

- $\bullet \ \emptyset$  Empty Set
- $\bullet~\mathbb{N}$  Set of nonnegative integers  $\{0,1,2\ldots\}$
- $\bullet~\mathbb{Z}$  Set of integers  $\{-2,-1,0,1,2\ldots\}$
- $\mathbb{Q}$  Set of rational numbers that can be expressed as p/q where  $p, q \in \mathbb{Z}$  and  $q \neq 0$ .  $\{-10.1, -1.2, 0, 5.5, 15...\}$
- $\mathbb{R}$  Set of real numbers  $\{e, \pi, \sqrt{2}, 2, 5.4\}$
- $\mathbb{C}$  Set of complex numbers  $\{2+5i, -i, 1, 23.3, \sqrt{2}\}$

**Comparing sets**: A is a subset of B ( $A \subseteq B$ ) iff every element of A is an element of B. E.g.  $A = \{a, b\}$  and  $B = \{a, b, c\}$ , then  $A \subseteq B$ . Every set is a subset of itself i.e.  $A \subseteq A$ .

A is a proper subset of B ( $A \subset B$ ) iff A is a subset of B, and A is not equal to B,

 $\begin{array}{l} {\sf Q}: \mbox{ Is } \{1,4,2\} \subset \{2,4,1\}. \mbox{ Is } \{1,4,2\} \subseteq \{2,4,1\} \mbox{ Ans: No, Yes } \\ {\sf Q}: \mbox{ Is } \mathbb{N} \subset \mathbb{Z}? \mbox{ Is } \mathbb{C} \subset \mathbb{R}? \mbox{ Ans: Yes, No} \\ {\sf Q}: \mbox{ What is } |\emptyset|? \mbox{ Ans: } 0 \end{array}$ 

**Union**: The union of sets A and B consists of elements appearing in A OR B. If  $A = \{1, 2, 3\}$  and  $B = \{3, 4, 5\}$ , then  $A \cup B = \{1, 2, 3, 4, 5\}$ .

**Intersection**: The intersection of sets A and B consists of elements that appear in both A AND B. If  $A = \{1, 2, 3\}$  and  $B = \{3, 4, 5\}$ , then  $A \cap B = \{3\}$ .

**Set difference**: The set difference of A and B consists of all elements that are in A, but not in B.  $A = \{1, 2, 3\}$  and  $B = \{3, 4, 5\}$ , then  $A \setminus B = A - B = \{1, 2\}$ .  $B \setminus A = B - A = \{4, 5\}$ .

**Complement**: Given a domain (or universe) D such that  $A \subset D$ , the complement of A consists of all elements that are not in A.  $D = \mathbb{N}$ ,  $A = \{1, 2, 3\}$ .  $A \subset D$  and  $\overline{A} = \{0, 4, 5, 6, \ldots\}$ .

 $A \cup \overline{A} = D$ ,  $A \cap \overline{A} = \emptyset$ ,  $A \setminus \overline{A} = A$ .

Q:  $D = \mathbb{N}$ ,  $A = \{1, 2, 3\}$  and  $B = \{3, 4, 5\}$ . Compute  $\overline{A \cap B}$ ,  $(B \setminus A) \cup (A \setminus B)$ .

Ans:  $\overline{A \cap B} = \{0, 1, 2, 4, 5, ...\}, (B \setminus A) \cup (A \setminus B) = \{1, 2, 4, 5\}$ 

**Disjoint sets**: Two sets are *disjoint* iff  $A \cap B = \emptyset$ .

**Symmetric Difference**:  $A \Delta B$  is the set that contains those elements that are either in A or in B, but not in both.

Q: Show  $A\Delta B$  on a Venn diagram. For  $A = \{1, 2, 3\}$  and  $B = \{3, 4, 5\}$ , compute  $A\Delta B$ .

**Cartesian product**: of sets is a set consisting of ordered pairs (*tuples*), i.e.  $A \times B = \{(a, b) \text{ s.t. } a \in A, b \in B\}.$ 

 $A = \{1, 2, 3\} \text{ and } B = \{3, 4, 5\}.$  $A \times B = \{(1, 3), (1, 4), (1, 5), (2, 3), (2, 4), (2, 5), (3, 3), (3, 4), (3, 5)\}.$ 

Similarly,  $A \times B \times C = \{(a, b, c) \text{ s.t. } a \in A, b \in B, c \in C\}.$ 

**Q**. Is  $A \times B = B \times A$ ? Ans: No. The order matters

#### **Distributive Law**: $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$

 $z \in A \cap (B \cup C)$ iff  $z \in A$  AND  $z \in (B \cup C)$ iff  $z \in A$  AND  $(z \in B \text{ OR } z \in C)$ Use the distributivity of AND over OR, x AND (y OR z) = (x AND y) OR (x AND z).

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iff (z \in A \text{ AND } z \in B) \text{ OR } (z \in A \text{ AND } z \in C)
iff z \in (A \cap B) \text{ OR } z \in (A \cap C)
iff z \in (A \cap B) \cup (A \cap C)
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# Questions?

A function assigns an element of one set, called the *domain*, to an element of another set, called the codomain s.t. for every element in the domain, there is exactly one element in the codomain. If A, B are sets then the function  $f : A \to B$ . Here A is the domain and B is the codomain. If  $a \in A$ , and  $b \in B$ , and f(a) = b, we say the function f maps a to b, b is the value of f at argument a, f assigns the element b to a, b is the image of a, a is the preimage of b.  $A = \{a, b, c, \dots z\}, B = \{1, 2, 3, \dots 26\}$ , then we can define a function  $f : A \rightarrow B$  such that f(a) = 1, f(b) = 2. f thus assigns a number to each letter in the alphabet. Consider  $f : \mathbb{R} \to \mathbb{R}$  s.t. for  $x \in \mathbb{R}$ ,  $f(x) = x^2$ .  $f(2.5) = 6.25 \in \mathbb{R}$ .

A function cannot assign different elements in the codomain to the same element in the domain. For example, if f(a) = 1 and f(a) = 2, the f is not a function.

A function that assigns a value to every element in the domain is called a *proper* function, while one that does not necessarily do so is called a *partial* function.

For  $x \in \mathbb{R}$ ,  $f(x) = 1/x^2$  is a partial function because no value is assigned to x = 0, since 1/0 is undefined.

Q: For  $x \in \mathbb{R}_+$ , consider  $f(x) = \sqrt{x}$ . Is f a function? Ans: Yes Q: For  $x \in [-1, 1], y \in \mathbb{R}$ , consider g(x) = y s.t.  $x^2 + y^2 = 1$ . Is g a function? Ans: No Q: For  $x \in \{-1, 1\}, y \in \mathbb{R}$ , consider g(x) = y s.t.  $x^2 + y^2 = 1$ . Is g a function? Ans: Yes

### **Functions**

Can define a function with a set as the argument. For a set  $S \in D$ ,  $f(S) := \{x | \forall s \in S, x = f(s)\}.$ 

 $A = \{a, b, c, \dots z\}, B = \{1, 2, 3, \dots 26\}. f : A \rightarrow B$  such that  $f(a) = 1, f(b) = 2, \dots$  $f(\{e, f, z\}) = \{5, 6, 26\}.$ 

If D is the domain of f, then range(f) := f(D) = f(domain(f)).

Q: If  $f : \mathbb{N} \to \mathbb{R}$ , and  $f(x) = x^2$ . What is the domain and codomain of f? What is the range? Ans:  $\mathbb{N}$ ,  $\mathbb{R}$ ,  $\{0, 1, 4, 9, ...\}$ 

Q: Consider  $f : \{0,1\}^5 \to \mathbb{N}$  s.t. f(x) counts the length of a left to right search of the bits in the binary string x until a 1 appears. f(01000) = 2.

What is f(00001), f(00000)? Is f a total function? Ans: 5, undefined, No.

**Surjective functions**: If  $f : A \to B$  is a surjective function, then for every  $b \in B$ , there exists an  $a \in A$  s.t. f(a) = b.

For surjective functions,  $|\#arrows| \ge |B|$ .

Since each element of A is assigned at most one value, and some need not be assigned a value at all,  $|\#arrows| \le |A|$ .

Hence, if f is a surjective function, then  $|A| \ge |B|$ .

 $A = \{a, b, c, \ldots z, \alpha, \beta, \gamma, \ldots\}, B = \{1, 2, 3, \ldots 26\}.$   $f : A \to B$  such that f(a) = 1,  $f(b) = 2, \ldots, f$  does not assign any value to the Greek letters. For every number in B, there is a letter in A. Hence, f is surjective. And |A| > |B|. **Injective functions**: If f is an injective function, then  $\forall a \in A$ , there is a *unique*  $b \in B$  s.t. f(a) = b.

Hence, |#arrows $| = |A| \le |B|$ . Hence, if f is a injective function, then  $|A| \le |B|$ .

 $A = \{a, b, c, \dots z\}, B = \{1, 2, 3, \dots 26, 27, \dots 100\}.$   $f : A \to B$  such that f(a) = 1,

f(b) = 2, ... No element in A is assigned values 27, 28, ..., and for every letter in A, there is a number in B. Hence, f is injective. And |A| < |B|.

**Bijective functions**: If f is a bijective function, then it is both surjective and injective, implying that |A| = |B|.

 $A = \{a, b, c, \dots z\}, B = \{1, 2, 3, \dots 26\}.$   $f : A \to B$  such that  $f(a) = 1, f(b) = 2, \dots$  Every element in A is assigned a value in B and for every element in B, there is a value in A that is mapped to it. f is bijective. And |A| = |B|.

Converse of the previous statements is also true.

- If  $|A| \ge |B|$ , then it's always possible to define a surjective function  $f : A \to B$ .
- If  $|A| \leq |B|$ , then it's always possible to define a injective function  $f : A \rightarrow B$ .
- If |A| = |B|, then it's always possible to define a bijective function  $f : A \rightarrow B$ .

Q: Recall that the Cartesian product of two sets  $S = \{s_1, s_2, \ldots, s_m\}$ ,  $T = \{t_1, t_2, \ldots, t_n\}$  is  $S \times T := \{(s, t) | s \in S, t \in T\}$ . Construct a bijective function  $f : (S \times T) \rightarrow \{1, \ldots, nm\}$ , and prove that  $|S \times T| = nm$ .

Ans:  $f(s_1, t_1) = 1$ ,  $f(s_1, t_n) = n$ ,  $f(s_2, t_1) = n + 1$ , and so on.  $f(s_i, t_j) = n(i-1) + j$ . Since f is bijective,  $|S \times T| = |\{1, ..., nm\}| = nm$ .

# Questions?

**Examples**: (a, b, a), (1,3,4), (4,3,1)

An element can appear twice. E.g.  $(a, a, b) \neq (a, b)$ .

The order of the elements does matter. E.g.  $(a, b) \neq (b, a)$ .

Q: What is the size of (1, 2, 2, 3)? What is the size of  $\{1, 2, 2, 3\}$ ? Ans: 4, 3.

**Sets and Sequences**: The Cartesian product of sets  $S \times T \times U$  is a set consisting of all sequences where the first component is drawn from *S*, the second component is drawn from *T* and the third from *U*.  $S \times T \times U = \{(s, t, u) | s \in S, t \in T, u \in U\}$ .

Q: For set  $S = \{0, 1\}$ ,  $S^3 = S \times S \times S$ . Enumerate S. What is  $|S^3|$ ?

Ans:  $S^3 = \{(0,0,0), (0,0,1) \dots (1,1,1)\}, |S^3| = 8$ 

### **Counting Sets - Example**

Suppose we want to buy 10 donuts. There are 5 donut varieties – chocolate, lemon-filled, sugar, glazed, plain. Let A be the set of ways to select the 10 donuts. Each element of A is a potential selection. For example, 4 chocolate, 3 lemon, 0 sugar, 2 glazed and 1 plain.

Let's map each way to a string as follows: 0000 000

chocolate lemon sugar glazed plain

Lets fix the ordering – chocolate, lemon, sugar, glazed and plain, and abstract this out further to get the sequence: 00001000110010.

Hence, each way of choosing donuts is mapped to a binary sequence of length 14 with exactly 4 ones. Now, let B be all 14-bit sequences with exactly 4 ones. An element of B is 11110000000000.

Q: The above sequence corresponds to what donut order? Ans: All plain donuts.

For every way to select donuts, we have an equivalent sequence in B. And every sequence in B implies a unique way to select donuts. Hence, the above mapping from  $A \rightarrow B$  is a bijective function.

### Counting Sets - using a bijection

Hence, |A| = |B|, meaning that we have reduced the problem of counting the number of ways to select donuts to counting the number of 14-bit sequences with exactly 4 ones.

**General result**: The number of ways to choose *n* elements with *k* available varieties is equal to the number of n + k - 1-bit binary sequences with exactly k - 1 ones.

Q: There are 2 donut varieties – chocolate and lemon-filled. Suppose we want to buy only 2 donuts. Use the above result to count the number of ways in which we can select the donuts? What are these ways?

Ans: Since n = 2, k = 2, we want to count the sequences with exactly 1 one in 3-bit sequences.  $\{(0, 0, 1), (1, 0, 0), (0, 1, 0)\}$ .

Q: In the above example, I want at least one chocolate donut. What are the types of acceptable 3-bit sequences with this criterion? How many ways can we do this?

Ans: We want to count the number of 3-bit sequences that start with zero and have exactly 1 one in them. So  $\{(0,1,0), (0,0,1)\}$ .

Suppose the university offers Math courses (denoted by the set M), CS courses (set C) and Statistics courses (set S). We need to pick one course from each subject – Math, CS and Statistics. What is the number of ways we can select we can select the 3 courses?

The above problem is equivalent to counting the number of sequences of the form (m, c, s) that maps to choose the Math course m, CS course c and Stats course s.

Recall that the product of sets  $M \times C \times S$  is a set consisting of all sequences where the first component is drawn from M, the second component is drawn from C and the third from S.  $M \times C \times S = \{(m, c, s) | m \in M, c \in C, s \in S\}$ . Hence, counting the number of sequences is equivalent to computing  $|M \times C \times S|$ .

**Product Rule**:  $|M \times C \times S| = |M| \times |C| \times |S|$ .

Using the above equivalence, the number of sequences and hence, the number of ways to select the 3 courses is  $|M| \times |C| \times |S|$ .

What is the number of length *n*-passwords that can be generated if each character in the password is allowed to be lower-case letter?

Each possible sequence is of the form (a, b, d, ...,) where the first element in the sequence can be selected from the  $\{a, b, ..., z\}$  set. Similar reasoning holds for each element.

Using the equivalence between sequences and products of sets, counting the number of such sequences is equivalent to computing  $|\{a, b, \dots z\} \times \{a, b, \dots z\} \times \{a, b, \dots z\} \dots |$ .

Using the product rule,  $|\{a, b, \dots z\} \times \{a, b, \dots z\} \times \{a, b, \dots z\} \dots | = |\{a, b, \dots z\}| \times |\{a, b, \dots z\}| \times \dots = 26^n$ .

Let R be the set of rainy days, S be the set of snowy days and H be the set of really hot days in 2022. A bad day can be either rainy, snowy or really hot. What is the number of good days?

Let B be the set of bad days.  $B = R \cup S \cup H$ , and we want to estimate  $|\overline{B}|$ . |D| = 365.  $|\overline{B}| = |D| - |B| = 365 - |B| = 365 - |R \cup S \cup H|$ .

Since the sets R, S and H are disjoint,  $|R \cup S \cup H| = |R| + |S| + |H|$ , and hence the number of good days = 365 - |R| - |S| - |H|.

**Sum rule**: If  $A_1, A_2 \dots A_m$  are disjoint sets, then,  $|A_1 \cup A_2 \cup \dots \cup A_m| = \sum_{i=1}^m |A_i|$ .

What is the number of passwords that can be generated if the (i) first character is only allowed to be a lower-case letter, (ii) each subsequent character in the password is allowed to be lower-case letter or digit (0 - 9) and (iii) the length of the password is required to be between 6-8 characters?

Let  $L = \{a, b, ..., z\}$  and  $D = \{0, 1, 2, ...\}$ . Using the equivalence between sequences and products of sets, the set of passwords of length 6 is given by  $P_6 = L \times (L \cup D)^5$ . Using the product rule,  $|P_6| = |L| \times (|L \cup D|)^5 = |L| \times (|L| + |D|)^5$ .

Since the total set of passwords are  $P = P_6 \cup P_7 \cup P_8$ ,  $|P| = |P_6| + |P_7| + |P_8| = |L| \times [(|L| + |D|)^5 (1 + (|L| + |D|) + (|L| + |D|)^2)] = 26 \times 36^5 \times [1 + 36 + 1296].$ 

# Questions?